

Propositional Logic

A *proposition* is a statement that is T or F.

Boolean variables p, q, r, \dots denote propositions.

Logical operators form *compound propositions*.

Truth tables evaluate compound propositions.

| | negation | conjunction | disjunction |
|-------------|----------|--------------|-------------|
| $p \quad q$ | $\neg p$ | $p \wedge q$ | $p \vee q$ |
| $T \quad T$ | F | T | T |
| $T \quad F$ | F | F | T |
| $F \quad T$ | T | F | T |
| $F \quad F$ | T | F | F |

| | exclusive or | implication | biconditional |
|-------------|--------------|-------------------|-----------------------|
| $p \quad q$ | $p \oplus q$ | $p \rightarrow q$ | $p \leftrightarrow q$ |
| $T \quad T$ | F | T | T |
| $T \quad F$ | T | F | F |
| $F \quad T$ | T | T | F |
| $F \quad F$ | F | T | T |

Propositional Equivalences

A compound proposition is:

a *tautology* if it is always true,

a *contradiction* if it is always false, or

a *contingency*, otherwise.

| p | q | $p \rightarrow q$ | $q \rightarrow (p \rightarrow q)$ |
|-----|-----|-------------------|-----------------------------------|
| T | T | T | T |
| T | F | F | T |
| F | T | T | T |
| F | F | T | T |

| p | q | $p \wedge q$ | $\neg p$ | $(p \wedge q) \rightarrow \neg p$ |
|-----|-----|--------------|----------|-----------------------------------|
| T | T | T | F | F |
| T | F | F | F | T |
| F | T | F | T | T |
| F | F | F | T | T |

| p | q | $p \oplus q$ | $p \leftrightarrow q$ | $(p \oplus q) \wedge (p \leftrightarrow q)$ |
|-----|-----|--------------|-----------------------|---|
| T | T | F | T | F |
| T | F | T | F | F |
| F | T | T | F | F |
| F | F | F | T | F |

p is *logically equivalent* to q if $p \leftrightarrow q$ is a tautology.

$p \leftrightarrow q$ or $p \equiv q$ denotes logical equivalence.

Use truth tables to determine logical equivalence.

$$p \rightarrow q \equiv \neg p \vee q$$

| p | q | $\neg p$ | $\neg p \vee q$ | $p \rightarrow q$ |
|-----|-----|----------|-----------------|-------------------|
| T | T | F | T | T |
| T | F | F | F | F |
| F | T | T | T | T |
| F | F | T | T | T |

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

| p | q | $p \vee q$ | $\neg(p \vee q)$ | $\neg p$ | $\neg q$ | $\neg p \wedge \neg q$ |
|-----|-----|------------|------------------|----------|----------|------------------------|
| T | T | T | F | F | F | F |
| T | F | T | F | F | T | F |
| F | T | T | F | T | F | F |
| F | F | F | T | T | T | T |

This equivalence is one of De Morgan's laws.

The other one is $\neg(p \wedge q) \equiv \neg p \vee \neg q$