Chapter 2: ARCHITECTURES

Software architectures and System architectures

Logical organization and Physical realization



Thanks to the authors of the textbook [**TS**] for providing the base slides. I made several changes/additions. These slides may incorporate materials kindly provided by Prof. Dakai Zhu.

So I would like to thank him, too.

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Chapter 2: ARCHITECTURES

- ARCHITECTURAL STYLES (SOFTWARE ARCHITECTURES)
- SYSTEM ARCHITECTURES
 - Centralized Architectures
 - Decentralized Architectures
 - Hybrid Architectures

ARCHITECTURES VERSUS MIDDLEWARE

- Interceptors
- General Approaches to Adaptive Software

SELF-MANAGEMENT IN DISTRIBUTED SYSTEMS

- The Feedback Control Model
- Example: Systems Monitoring with Astrolabe
- Example: Differentiating Replication Strategies in Globule

Objectives

- To learn how to organize a distributed system whose components are dispersed across multiple machines
- To understand the differences between
 - software architecture (logical organization) and
 - system architecture (physical realization)
- To understand trade-offs when providing distribution transparency
- To understand adaptability and self-mng issues and mechanisms for flexibility and efficiency

Software Architecture

(Architectural Style, Logical organization)

- Divide the system into logically different software components, distribute them over multiple machines, and allow them to communicate through connectors
 - Component: a modular unit with well-defined required and provided interfaces,
 - Connector: a mechanism that mediates communication, coordination, and cooperation (e.g., RPC, msg passing)
- Using components and connectors, we can create different configurations, which are classified into the following architectural styles:
 - Layered
 - Object-based
 - Event-based
 - Data-centered

All try to achieve distributed transparency at a reasonable level and

Each style would be more appropriate for a different application

Software Architecture

(Architectural Style, Logical organization)

Layered style

 used for client-server systems, request/reply model

Object-based

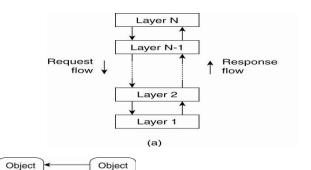
 Used for distributed object systems, request/reply model

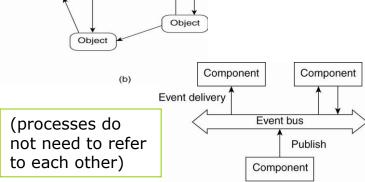
Event-based:

- Publish/subscribe systems
- Loosely coupled components
 - decoupled in space or referentially decoupled

Data-centered:

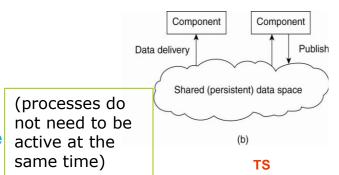
- Communicate through common repository (e.g., shared distributed file system)
- Can be combined with event-based, yielding shared dataspace
 - processes are now decoupled in space and time





Method call

Object



(a)

System Architecture (Physical realization)

- Consider how and where to place software components and realize their interactions
- There are three major physical realization approaches:

Centralized

client-server

Decentralized

P2P (Structured vs. unstructured)

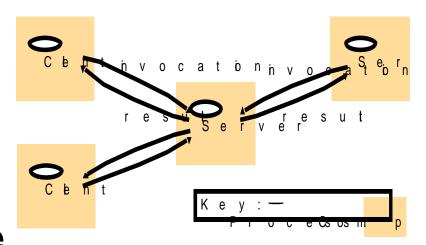
Hybrid:

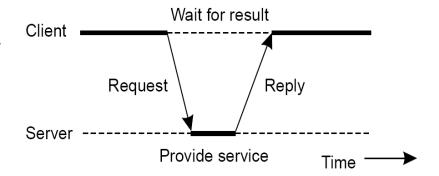
combination of centralized and P2P

System Architecture: Client-Server

- There are processes offering services (servers)
- There are processes that use services (clients)
- Clients and servers can be on different machines
- Clients follow request/reply model to use services
- Connection-oriented vs. connectionless

(most use TCP vs UDP)

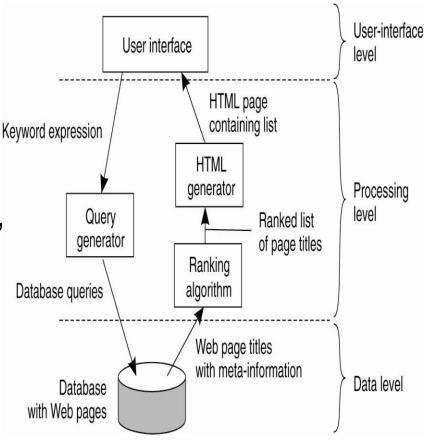




Application Layering (logical)

How to draw a clear line between client end server?

- User-interface layer
 - units for an application's user interface
- Processing layer:
 - functions of an application,
 i.e. without specific data
- Data layer:
 - data that a client wants to manipulate through the application components



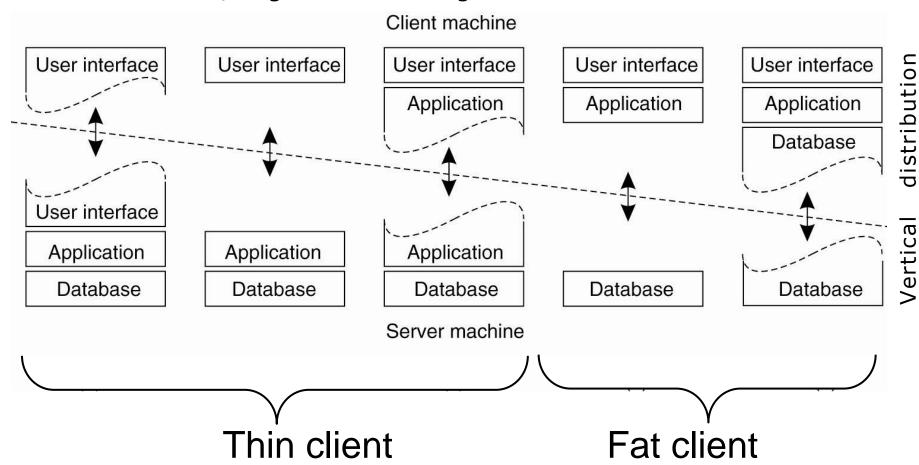
Observation: layering is found in many distributed information systems, using traditional database technology and accompanying applications.

Multitiered Architectures (physical realization)

How to place the three layers on client and server?

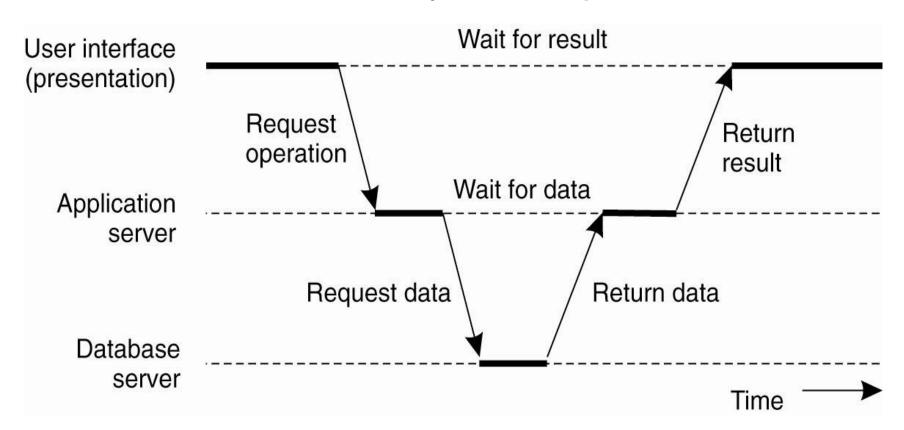
Single-tiered: dumb terminal/mainframe configuration

Two-tiered: client/single-server configuration



Multitiered Architectures (physical realization)

- The server part could be distributed over multiple machines,
- Three-tiered: each layer on separate machine

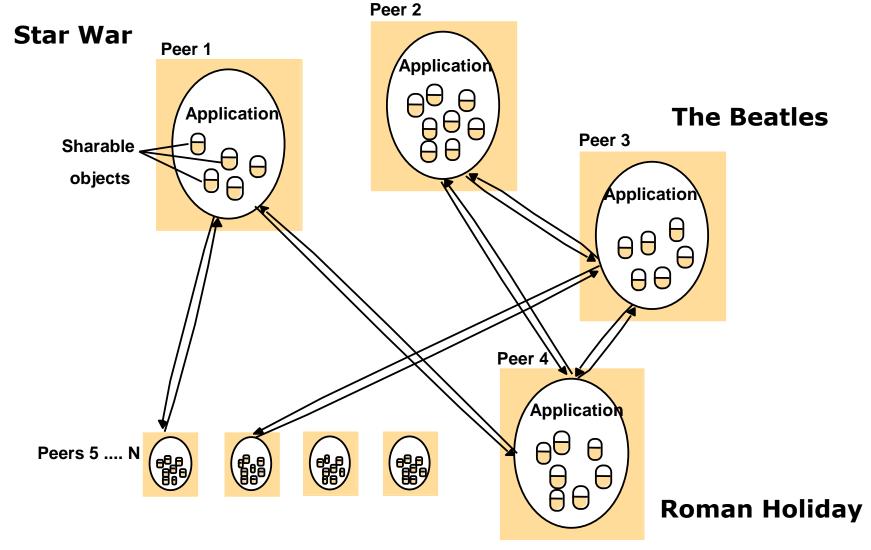


Decentralized Architectures: P2P Systems

- Multitiered architectures do vertical distribution:
- We place logically different components in client-server (i.e., user interface, processing, data) on different machines
- Processes are not equal and Interactions are asymmetric
 - One acts as client while the other acts as server
- Traditional approach

- P2P architectures do horizontal distribution:
- We split up clients and servers into logically equivalent parts and let each part operate on its own share
- Processes are equal and Interactions are symmetric
 - Each acts as both client and server
- Tremendous growth in the last couple of years

Decentralized Architectures: P2P Systems



Decentralized Architectures: P2P Systems

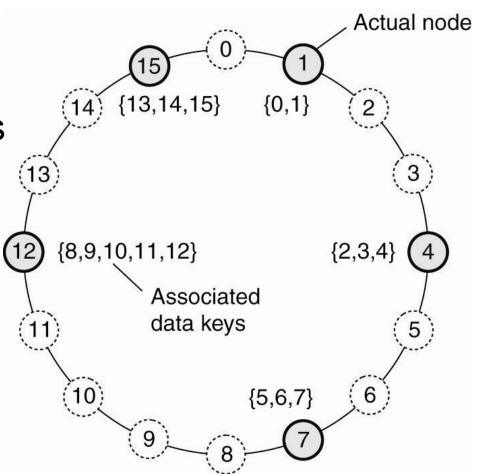
- Given the symmetric behavior, the key question is how to organize processes in an overlay network, where links are usually TCP channels...
- How about fully connected overlay network? -/+
- There are three approaches to organize nodes into overlay networks through which data is routed
 - Structured P2P: nodes are organized following a specific distributed data structure and deterministic algorithms
 - Unstructured P2P: randomly selected neighbors
 - Hybrid P2P: some nodes are appointed special functions in a well-organized fashion

Structured P2P Systems

- Distributed Hash Table (DHT) is the most used one
 - Assume we have a large ID space Ω (e.g., 128-bit)
 - Assign random keys from Ω to data items
 - Assign random identifiers from Ω to nodes
 - The crux of every DHT is to implement an efficient and deterministic scheme that maps the key of a data item to node ID
 - When looking up a data item, the system should route the request to the associated node and return the network address of that node
- Example: Chord

A DHT Example: Chord

- Chord organizes the nodes in a structured overlay network such as a logical ring, and data item with key k is mapped to a node with the smallest ID >= k.
- This node is called as the successor of key k and denoted by succ(k)



LOOKUP(key=8)?

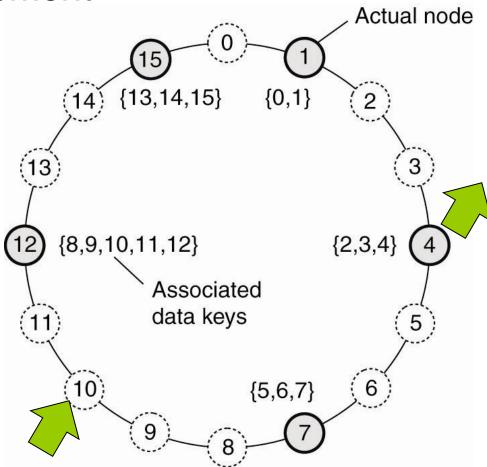
This should return *succ(8)* which is node 12. (Details of how this is done is in Ch 5)

A DHT Example: Chord

Membership management

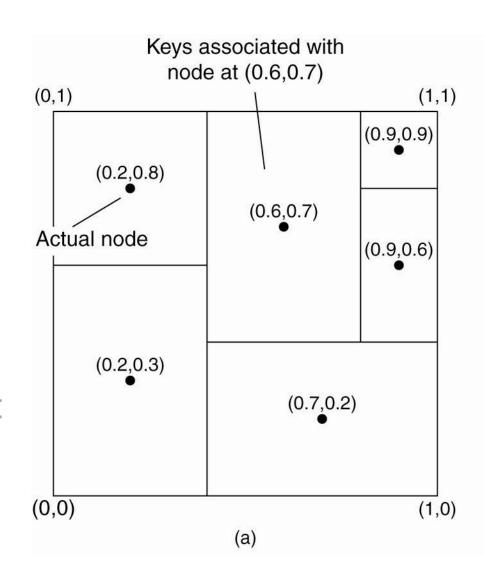
- Join
- Leave
- Lookup(key)

(search and routing is in Ch 5)



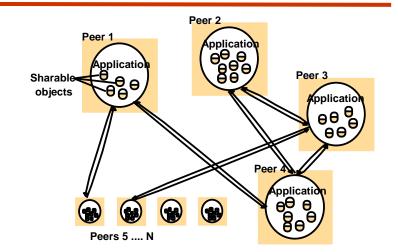
Another DHT Example: CAN

- Content Addressable Network
- Organize nodes in a ddimensional space and let every node take the responsibility for data in a specific region.
- When a node joins, split a region.
- When a node leaves, merge regions.



Unstructured P2P Systems

- Maintain a random graph
- Data items are randomly placed on nodes
- How to do Lookup?
 - flooding
- Membership management
 - Join
 - Get a random list (from a well-known list or server)
 - Contact these nodes and run the algorithm presented next
 - Leave
 - Easy just leave...



How to maintain random graph

- Let each peer maintain a partial view of the network, consisting of c other nodes
- Each node P periodically selects a node Q from its partial view
- P and Q exchange information and exchange members from their respective partial views

Actions by active thread (periodically repeated):

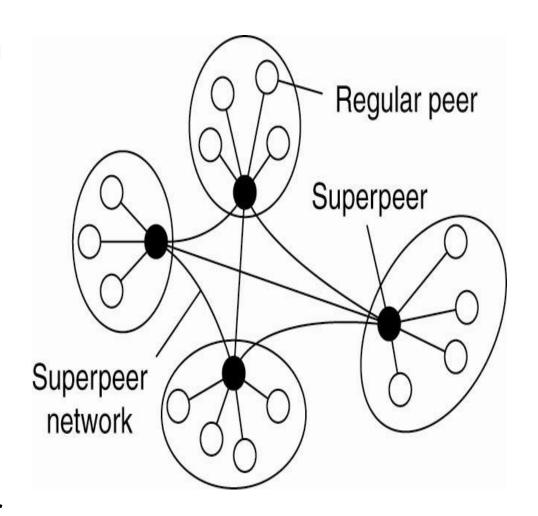
```
select a peer P from the current partial view;
if PUSH_MODE {
    mybuffer = [(MyAddress, 0)];
    permute partial view;
    move H oldest entries to the end;
    append first c/2 entries to mybuffer;
    send mybuffer to P;
} else {
    send trigger to P;
}
if PULL_MODE {
    receive P's buffer;
}
construct a new partial view from the current one and P's buffer;
increment the age of every entry in the new partial view;
```

Actions by passive thread:

```
receive buffer from any process Q;
if PULL_MODE {
    mybuffer = [(MyAddress, 0)];
    permute partial view;
    move H oldest entries to the end;
    append first c/2 entries to mybuffer;
    send mybuffer to P;
}
construct a new partial view from the current one and P's buffer;
increment the age of every entry in the new partial view;
    (b)
```

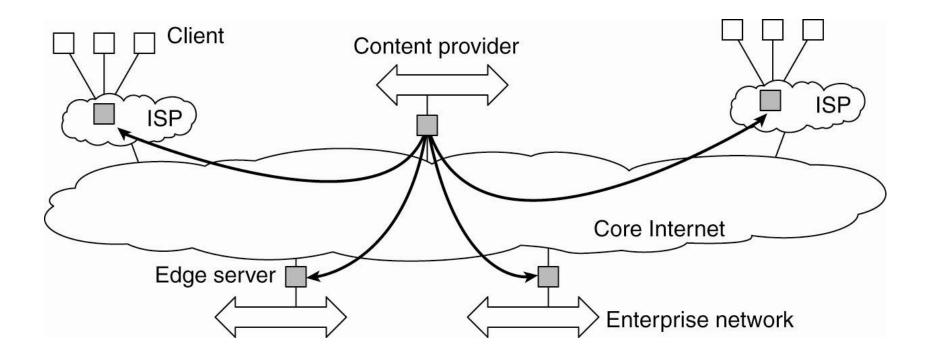
Superpeers

- When random graph gets bigger, it will be very hard to perform look up
- Use supperpeers to maintain an index
- Join/leave is easy
- How about Lookup?
- Regular peers may elect the supperpeer (Ch 6)



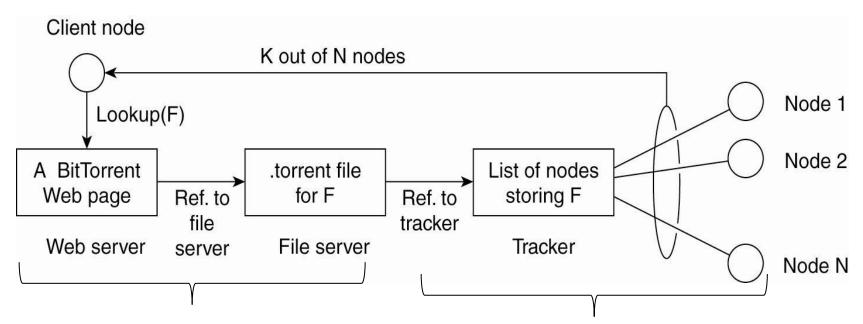
Hybrid Architectures: Edge-server systems

- Content Distribution Network (CDN)
- Edge servers can be used to optimize content distribution



Hybrid Architectures: Collaborative Distributed Systems

- Combining a P2P with a client-server architecture
- Basic idea: a node identifies where to download a file from and joins a swarm of downloaders; who get file chunks in parallel from the source, and distribute these chunks amongst each other



WHERE MIDDLEWARE FITS IN ALL THESE ARCHITECTURES?

Architectures Vs. Middleware

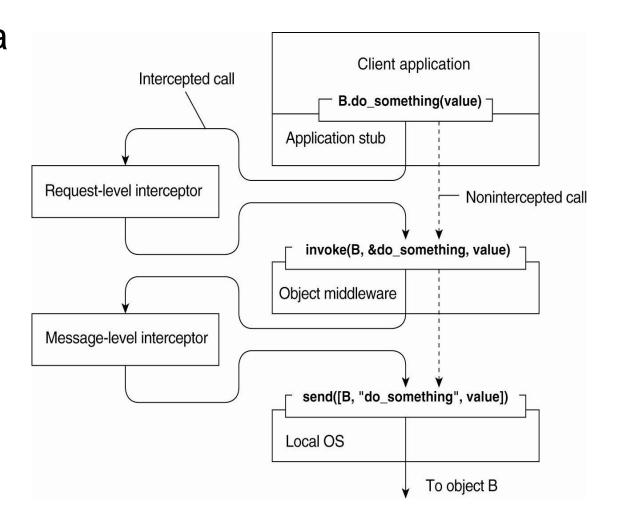
- Middleware is between application and local OS and provides some degree of transparency
- In practice, middleware systems follow a specific architectural style (software architecture, logical organization):
 - Layered
 - Object-based
 - Data centered
 - Event-based

Adv/DisAdv

- + makes app design simple
- may not be optimized for what an app needs
- adding more features complicates the middleware
 - CORBA was initially object-based, later added msg passing
- Middleware should be adaptable to applications
 - Several different versions, configurable, separate policy and mechanisms

How to achieve adaptability?

Interceptors: a software construct that will break the usual flow of control and allow other (app specific) code to be executed



General Approaches to Adaptive Software

In many cases, distributed systems/applications are developed according to a specific architectural style. The chosen style may not be optimal in all cases Then there is a need to (dynamically) adapt the behavior of the middleware.

Three basic approaches to adaptive software:

Separation of concerns:

Try to separate extra functionalities and later glue them together into a single implementation → aspect-oriented SW ,only toy examples so far.

Computational reflection:

Let a program inspect itself at runtime and adapt/change its settings dynamically if necessary → mostly at language level and applicability unclear.

Component-based design:

- Organize a distributed application through components that can be dynamically replaced when needed (complex for DS, components are not independent)
- Do we really need adaptive software or adaptive system that reacts to changes (self-management)

Do we really need adaptive software?

- Software should expect all the environment changes and should have code in it to handle them
- DS should be able to react to changes in environment by switching policies or mechanisms in the system
- The challenge is how to achieve this reactive behavior without human intervention

Self-managing Distributed Systems

- Distinction between system and software architectures blurs when automatic adaptivity needs to be taken into account:
- Self-configuration
- Self-managing
- Self-healing
- Self-optimizing
- Self-*

Warning

There is a lot of hype going on in this field of autonomic computing.

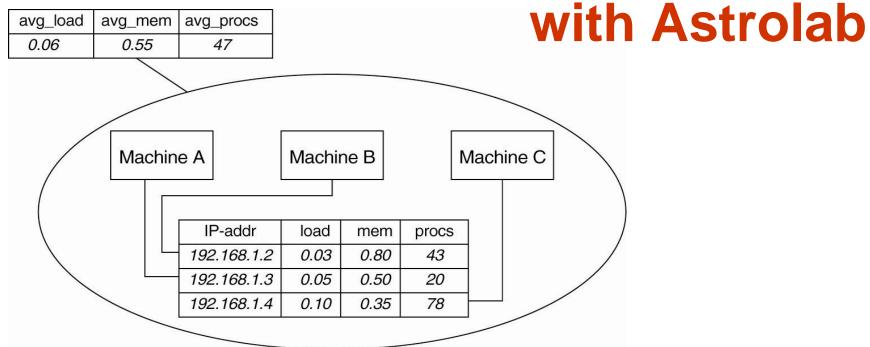
Feedback Control Model

Logical view! Logical realization Physical realization might be distributed might be In many cases, self-* systems use a feedback control loop. Uncontrollable parameters (disturbance / noise) Initial configuration Corrections Observed output Core of distributed system +/-+/-Reference input Adjustment Metric estimation measures Analysis Measured output Adjustment triggers

if time permits

FEEDBACK CONTROL EXAMPLES

Example: Systems Monitoring



A general tool for observing system behavior

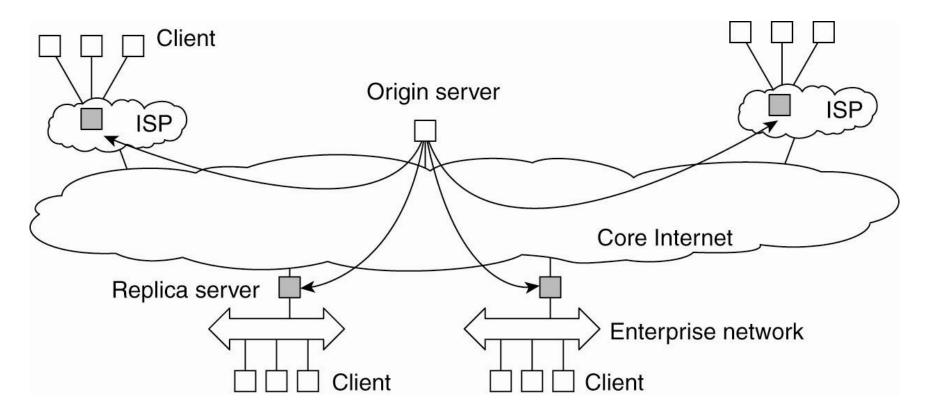
Organize hosts into a hierarchy of zones.

Collect information about each host and summarize it,

Exchange this information so all agents will see the same view.

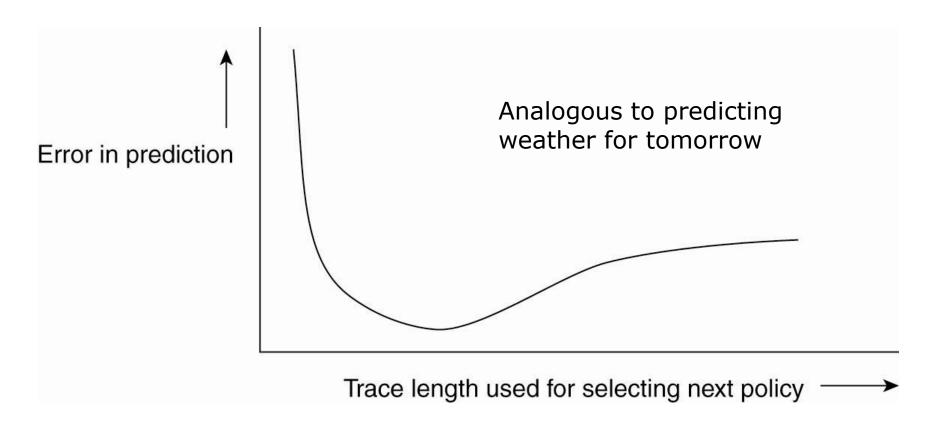
Example: Differentiating Replication Strategies in Globule (1)

A collaborative CDN tries to minimize performance by replicating web pages.



Example: Differentiating Replication Strategies in Globule (2)

■ Figure 2-19. The dependency between prediction accuracy and trace length.



Example: Automatic Component Repair Management in Jade

- Steps required in a repair procedure:
- Terminate every binding between a component on a nonfaulty node, and a component on the node that just failed.
- Request the node manager to start and add a new node to the domain.
- Configure the new node with exactly the same components as those on the crashed node.
- Re-establish all the bindings that were previously terminated.