

## Discrete Mathematical Structures CS 2233 Lecture Two

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## Course Introduction

- **Recall:** Homework one due Thursday 1/22
  - Section 1.1: 6d, 6e, 6g, 10, 26, 28a-d
  - Hand in at beginning of class. Work alone.
- Recitation will be held today, Thursday 1/15
  - I plan to use the Thursday and Tuesday Prior to a homework assignment being due to give you a chance to ask questions about it
  - If you find you want to attend a different recitation than the one you have registered for, that's OK

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## Syntax of Propositional Formulas

- Definition of *propositional formula*:
  - A propositional variable  $p$  is a propositional formula
  - The constants T and F are propositional formulas
  - If  $\phi$  and  $\psi$  are propositional formulas, then the following are also propositional formulas:
    - $(\phi)$
    - $\neg\phi$
    - $\phi \wedge \psi$
    - $\phi \vee \psi$
    - $\phi \rightarrow \psi$
    - $\phi \leftrightarrow \psi$
    - $\phi \oplus \psi$

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## Semantics of Propositional Formulas

- A formula defines a function from truth assignments to truth values
  - A truth assignment gives a truth value for each variable

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## Negation

- Def: Given a proposition  $p$ ,  $\neg p$  denotes the negation of  $p$ 
  - $\neg p$  means "it is not the case that  $p$ "
- Truth table for  $\neg p$ :

$p$	$\neg p$
T	F
F	T

↑  
Given truth assignments for  $p$

↑  
Resulting truth value of  $\neg p$

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## Conjunction

- Def: Given propositions  $p$  and  $q$ ,  $p \wedge q$  denotes the conjunction of  $p$  and  $q$ 
  - $p \wedge q$  means "p and q"
- Truth table for  $p \wedge q$ :

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

↑  
Given truth assignments

↑  
Resulting truth value

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## Disjunction

- Def: Given propositions  $p$  and  $q$ ,  $p \vee q$  denotes the disjunction of  $p$  and  $q$ 
  - $p \vee q$  means "p or q" (*inclusive or*)
- Truth table for  $p \vee q$ :

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Given truth assignments      Resulting truth value

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## Implication

- Def: Given propositions  $p$  and  $q$ ,  $p \rightarrow q$  is an implication
  - $p \rightarrow q$  means "p implies q"
  - $p$  is the *hypothesis, antecedent or premise*
  - $q$  is the *conclusion, consequence, or consequent*
  - Truth table for  $p \rightarrow q$ :

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

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## Understanding Implications

- Some readings of  $p \rightarrow q$ :
  - if  $p$ , then  $q$
  - $q$  if  $p$
  - $q$  when  $p$
  - $p$  only if  $q$
  - $q$  follows from  $p$
  - $q$  is a necessary condition for  $p$
  - $p$  is a sufficient condition for  $q$
  - $q$  is necessary for  $p$
  - $p$  is sufficient for  $q$

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## Formulas Related to $p \rightarrow q$

- Contrapositive:**  $\neg q \rightarrow \neg p$
- Converse:**  $q \rightarrow p$
- Inverse:**  $\neg p \rightarrow \neg q$
- Equivalent formula not using  $\rightarrow$ :  $\neg p \vee q$

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## Equivalence

- Two formulas are equivalent if they have the same semantics
  - They yield the same truth values on the same truth assignments

p	q	$\neg p$	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$	$\neg p \vee q$
T	T	F	F	T	T	T	T	T
T	F	F	T	F	F	T	T	F
F	T	T	F	T	T	F	F	T
F	F	T	T	T	T	T	T	T

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## Notation for Equivalence

- When  $\phi$  and  $\psi$  are we write  $\phi \equiv \psi$ 
  - e.g.,  $(\phi \rightarrow \psi) \equiv (\neg \phi \vee \psi)$
- Note that  $\equiv$  is not a logical connective (*i.e.*, a logical operator)
  - These are not formulas:
    - $\phi \equiv \psi$
    - $(\phi \equiv \psi) \wedge p$
- $\equiv$  is part of the *meta-language* we use for discussing formulas in the *object language* of propositional calculus

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## Exclusive Or

- Def: Given propositions  $p$  and  $q$ ,  $p \oplus q$  denotes the exclusive or of  $p$  and  $q$ 
  - $p \oplus q$  means "p or q, but not both"
- Truth table for  $p \oplus q$ :

$p$	$q$	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

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## Biconditionals

- Def: Given propositions  $p$  and  $q$ ,  $p \leftrightarrow q$  is a *biconditional*
  - $p \leftrightarrow q$  means "p if and only if q"
- Truth table for  $p \leftrightarrow q$ :

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

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## Translating English to Logic

- Fred can access the wireless network only if Fred has paid his tuition
  - Let  $a$  represent "Fred can access wireless"
  - Let  $t$  represent "Fred has paid his tuition"
  - $a \rightarrow t$

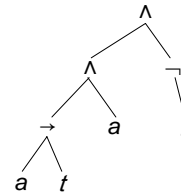
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## Expression Trees

- Example:  $((a \rightarrow t) \wedge a) \wedge \neg t$



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## A More Concise Truth Table

$((a \rightarrow t) \wedge a) \wedge \neg t$	$a$	$t$
T	T	T
T	F	T
F	T	F
F	F	F

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## A More Concise Truth Table

$((a \rightarrow t) \wedge a) \wedge \neg t$	$a$	$t$
T	T	T
T	F	T
F	T	F
F	F	F

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## A More Concise Truth Table

$((a \rightarrow t) \wedge a) \wedge \neg t$	$a$	$\neg t$	$t$
T	T	T	F
T	F	F	T
F	T	T	F
F	T	F	T

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## A More Concise Truth Table

$((a \rightarrow t) \wedge a) \wedge \neg t$	$a$	$\neg t$	$t$
T	T	T	F
T	F	F	T
F	T	T	F
F	T	F	T

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## Precedence

- Should  $\neg q \rightarrow \neg p$  be interpreted as
  - $(\neg q) \rightarrow (\neg p)$ , or as
  - $\neg(q \rightarrow (\neg p))$ ?
- Precedence gives rules for implicit parentheses

Op	Prec
$\neg$	1
$\wedge$	2
$\vee$	3
$\rightarrow$	4
$\leftrightarrow$	5

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## Consistency

- Intuitive requirement: specifications should not contain conflicting requirements
- Definition: A collection of propositional formulas is *consistent* if there is a truth assignment that makes each formula true

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## Example Inconsistent Spec

- Specification:
  - Fred can access the wireless network only if Fred has paid his tuition ( $a \rightarrow t$ )
  - Fred can access the wireless network ( $a$ )
  - Fred has not paid his tuition ( $\neg t$ )

$a$	$t$	$a \rightarrow t$	$\neg t$	$(a \rightarrow t) \wedge a \wedge \neg t$
T	T	T	F	F
T	F	F	T	F
F	T	T	F	F
F	F	T	T	F

- $(a \rightarrow t) \wedge a \wedge \neg t$  is a *contradiction*

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## Tautologies and Contradictions

- A compound proposition that is true for all truth assignments is a *tautology*
  - E.g.,  $(p \rightarrow q) \leftrightarrow (\neg q \rightarrow \neg p)$
- One that is false for all assignments is a *contradiction*

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