

# Discrete Mathematical Structures

## CS 3233 Lecture Fourteen

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# Business

- Homework assignment 3 is due tomorrow by 2 at my office
- Lecture schedule:
  - No lecture next week (9/26-30)
  - Meet with TA during lecture time ALL the following week (10/3-7)
- Questions??
- Outline for today
  - Composition of function and its inverse is identity
  - Cover two remaining slides from make up lecture
  - Revisit writing formulas stating that exactly one and exactly two elements of universe satisfy an arbitrary, given formula
  - Start section 2.1

# Function Inverses

- Theorem: Given  $f : A \xrightarrow[\text{onto}]{1-1} B$ , every  $b \in B$  has a unique pre-image  $a \in A$
- This justifies the following definition: Given

$$f : A \xrightarrow[\text{onto}]{1-1} B$$

- $f^{-1} : B \rightarrow A$  is the *inverse of*
- $f^{-1}(b) = a$  iff  $f(a) = b$

# Compositions

- Definition
  - Given  $g : A \rightarrow B$  and  $f : B \rightarrow C$ , the composition of  $f$  and  $g$ ,  $f \circ g : A \rightarrow C$ , is defined by  $f \circ g(a) = f(g(a))$
  - Note that no special properties of  $f$  and  $g$  are required for  $f \circ g$  to be defined.
    - For instance,  $f$  and  $g$  need not be injective, surjective, or bijective
    - However, if  $f$  and  $g$  have special properties, it often follows that  $f \circ g$  special properties as well
    - Study hint: think through these relationships

# Identity Function

- Given any set  $A$ , the identity function on  $A$ ,  $\iota : A \rightarrow A$ , is defined by  $\iota(a) = a$  for all  $a \in A$
- Theorem
  - Given any two sets  $A$  and  $B$  such that there exists a one-to-one correspondence  $f : A \rightarrow B$ ,  $f^{-1} \circ f : A \rightarrow A$  and  $f \circ f^{-1} : B \rightarrow B$  are the identity functions on  $A$  and  $B$  respectively

# Exercise 35, Section 3.2

- Theorem: If  $A$  is an uncountable set and  $A$  is a subset of  $B$ , then  $B$  is uncountable
- Proof:
  - Suppose  $B$  is countable
  - This means there is a sequence that enumerates  $B$
  - A sequence that enumerates  $A$  can be constructed by dropping the elements of  $B - A$ , yielding the desired contradiction

# Uncountability of the Reals

- Theorem
  - $\mathbb{R}$ , the set of real numbers, is uncountable
- Proof
  - Uses Georg Cantor's *diagonalization argument*
  - Outline
    - Assume for contradiction that the real interval  $[0,1]$  is enumerated by  $\{a_n\}$
    - Use  $\{a_n\}$  to construct a real in  $[0,1]$  that does not occur in  $\{a_n\}$
    - Idea: for each decimal place,  $n$ , in the representation of the constructed value, choose a decimal different from the  $n^{\text{th}}$  place of  $a_n$
    - The fact that the constructed value differs from each value  $a_n$  shows that  $\{a_n\}$  does not enumerate  $[0,1]$
    - Now use Exercise 35 to complete the proof that  $\mathbb{R}$  is uncountable

# Exactly One, Two

- Construct a formula that is true just in case  $p(x)$  holds for exactly one value of  $x$  in the universe of discourse
- Do same for two values holding

# Algorithms

- Read section 2.1
- Definition
  - An *algorithm* is a finite collection of precise instructions for performing a computation to solve a problem
  - Input and output values are elements of specified sets
- Desirable characteristics:
  - Definiteness: Steps are precisely defined
    - To be really precise, must use a formal *computational model*, such as a Turing machine or the lambda calculus
  - Effectiveness: It must be possible to perform each step using a bounded amount of time and storage space
    - Bounded means there is an amount of time that is always sufficient
  - Correctness, Termination, Generality

# Examples

- Finding maximum value in an input sequence of integers
- Searching: Determine whether a given value is contained in an input sequence of integers