

## Discrete Mathematical Structures CS 3233 Lecture Five

Prof. William Winsborough  
September 7, 2006

## Business

- Read Sections 1.6 and 1.7 by Tuesday (You can stop on p.97 before "Tilings".)
- **Recall: Homework Two:** due Tuesday September 12
  - Section 1.3: 10d, 10e, 14, 24c, 24d, 32a, 32b, 44
  - Section 1.4: 2a, 2c, 4, 10, 24, 46
  - Section 1.5: 4, 8
- Questions???
- Return Homework One
  - Will be discussed in recitation
- Practice quiz

7 September 2006

Winsborough CS 3233 Lecture 5

2

## Practice Quiz: In Which Numeric Domains Does each of the Following Hold?

- Domains
  - $\mathbb{Z}$  – The integers
  - $\mathbb{N}$  – The natural numbers (non-negative integers)
  - $\mathbb{R}^+$  – The positive reals
  - $\mathbb{R}^+ \cup \{0\}$  – The non-negative reals
- Formulas
  - $\forall x. \exists y. y < x$
  - $\exists x. \forall y. x \leq y$
  - $\forall x. \forall z. (x < z \rightarrow \exists y. (x < y) \wedge (y < z))$

7 September 2006

Winsborough CS 3233 Lecture 5

3

## Free and Bound Variables

- A variable occurrence in a formula is called *bound* if it lies in the scope of a quantification of that variable:
  - All bound:  $\forall x. \forall z. (x < z \rightarrow \exists y. (x < y) \wedge (y < z))$
  - Free  $y$ :  $\forall x. \forall z. (x < z \rightarrow (x < y) \wedge (y < z))$
  - Free  $x$  and bound  $x$ :  $x < 7 \wedge \forall x. (p(x) \rightarrow x > 1)$
- Otherwise a variable occurrence is called *free*

7 September 2006

Winsborough CS 3233 Lecture 5

4

## Open and Closed Formulas

- A formula is *closed* if all variable occurrences are bound
  - Otherwise, it is *open*
  - *E.g.*,  $\Phi(x) = \forall y. x \leq y$  is open
- When denoting a formula by a meta-variable like  $\Phi$ , we list the variables that occur free in the formula after the variable:  $\Phi(x)$ 
  - Can refer to an instance of  $\Phi(x)$  by providing values for free variables:  $\Phi(5) = \forall y. 5 \leq y$
  - (In what numeric domains does  $\Phi(0)$  hold?)

7 September 2006

Winsborough CS 3233 Lecture 5

5

## Theorems and Proofs

- A *theorem* is a statement (such as a formula) that can be shown to be true in all cases (a tautology)
- A proof is a demonstration that a statement is a theorem
- Example methods of proof
  - Construction of truth tables
  - Use of equivalences
    - By using these alone, can prove only logical equivalences
  - More general rules of inference

7 September 2006

Winsborough CS 3233 Lecture 5

6

## Important Related Terminology

- *Result*: often used to mean a theorem
- *Proposition*: a simple theorem, often presented without proof
- *Lemma*: a theorem whose main utility lies in helping to prove other, more interesting theorems
- *Corollary*: a theorem that follows easily from another more general theorem
- *Conjecture*: a statement that you suspect is true but that you do not yet have a proof for

7 September 2006

Winsborough CS 3233 Lecture 5

7

## Rules of Inference for Propositional Logic

- A general, systematic method of proving formulas
- See Table 1 p.66
  - Known equivalences can also be used in proofs
- Use rules of inference to show
  - These hypotheses:
    - If it does not rain or if it is not foggy, the sailing race will be held and the lifesaving demonstration will take place
    - If the race is held, the trophy will be awarded
    - The trophy was not awarded
  - Imply this conclusion:
    - It rained

7 September 2006

Winsborough CS 3233 Lecture 5

8

## Rules of Inference for Universally Quantified Statements

- Universal instantiation

$$\frac{\forall x. p(x)}{p(c)}$$

- For any c

- Universal generalization

$$\frac{p(c)}{\forall x. p(x)}$$

- Must show p(c) for arbitrary c

7 September 2006

Winsborough CS 3233 Lecture 5

9

## Rules of Inference for Existentially Quantified Statements

- Existential instantiation

$$\frac{\exists x. p(x)}{p(c)}$$

- c must be a new name (constant) that does not appear earlier in the proof

- Existential generalization

$$\frac{p(c)}{\exists x. p(x)}$$

- c can be any name

7 September 2006

Winsborough CS 3233 Lecture 5

10

## Universal Modus Ponens

- Combines propositional modus ponens with universal instantiation:

$$\frac{\forall x. \Phi(x) \rightarrow \Psi(x)}{\Phi(c) \quad \Psi(c)}$$

7 September 2006

Winsborough CS 3233 Lecture 5

11

## Using Rules of Inference for Quantified Statements

- Problems 14 and 15, p.73
- We won't always be so meticulous about identifying exactly which rules of inference are being used
  - But it's important to understand the fundamentals of how proofs can be constructed

7 September 2006

Winsborough CS 3233 Lecture 5

12