Modular Arithmetic

Arithmetic of remainders resulting from divisions of integers by a positive integer.

Definitionlet a and b be integers and m be apositive integer. a is congruent to b modulo mif m | (a-b).a = b (mod m)if a is congruent to b modulo ma
$$\neq$$
 b (mod m)if a is not congruent to b modulo ma \neq b (mod m)if a is not congruent to b modulo mExamples:is 2 = s (mod 3) ? bees 3 | (2-s)? yes.is 17 = 7 (mod s) ? bees 5 | (17-7)? yes.(Practice Problem 20.26 and 29)Section 4.1, problem 11: 12-hour clock. What is the time shown by the clock80 hours after 11:00? 80+11 (mod 12) = 7. The clock will show 7:00.40 hours before 12:00? 12:40 (mod 12) = -28 mod 12 = 36-28 (mod 12) = 8. The clock showed 8:00.Theorem 4.Let a and b be integers and m be a positive integer.a = b (mod m) if and only if a = b+km for someinteger k.proof: \Rightarrow Let a = b (mod m) \therefore m | (a-b) by def. \therefore (a-b) = k·m (by def-of divide notatim) k is an int \therefore a = b+km \leftarrow Let a = b+km \therefore (a-b) = k m or (a-b) is a multipk of m \therefore m | (a-b) \therefore a = b (mod m)

Theorem 3. Let a, b be integers, m be a positive integers.
a = b (mod m) if and only if (a mod m) = (b mod m)
=: Let a = b (mod m)
$$\Rightarrow$$
 a = b+km (from Theorem 4)
(1)
By Division Algorithm, there exist two integers y and y
o = $y < m$, such that
b = $q \cdot m + y$ as b mod m = y
 \therefore a = $(9m + x) + km$ [substitute for b in (i)]
a = $9m + km + y$
a = $(9r + k)m + x$ \therefore a mod m = y
 \therefore (a mod m) = (b mod m) = y .
 \Leftarrow Let (a mod m) = (b mod m)
Let a mod m = b mod m = x , o s $y < m$.
 \therefore a = $q_1 \cdot m + y$ $\&$ b = $q_2 \cdot m + y$
 $fry home inth, q_1 and y_2
 $a - b = (q_1 - q_2)m$
 \therefore m | a-5 \Rightarrow a = b (mod m)$

Theorem 5. Let
$$a, b, c, d$$
 be integers and m be a positive int:
If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then
(i) $a + c \equiv b + d \pmod{m}$
(ii) $a c \equiv b d \pmod{m}$
(ii) $a c \equiv b d \pmod{m} \Rightarrow a \equiv b + km \left(for some int k + Thm. 4\right)$
 $c \equiv d \pmod{m} \Rightarrow c \equiv d + l \cdot m \quad (for some int. 1)$
 $a + c \equiv b + d + (d + l m)$
 $a + c \equiv b + d + (k + l) m$
 $\therefore a + c \equiv b + d \pmod{m} \quad (Theorem 4)$
[Prove (iii]
Corollary 2 (i) $(a + b) \mod m \equiv (a \mod m + b \mod m) \mod m$
(i).
Fact: $a \equiv (a \mod m) \pmod{m}$. Why?
By the division algorithm, there exist two unique integers q and r such that

 $a = q m + r, 0 \le r < m.$

Since $m \mid a-r$, $a \equiv r \pmod{m}$, where $r = a \mod m$ (by definition)

a = (a mod m) (mod m) b = (b mod m) (mod m) :. atb = ((a mod m) + (b mod m)) mod m (Thm 5, i)

Problem 26

a=4(mod12) List 5 integers that are congruent to 4 modulo 12. 4, 16, 28, 40, 52 a = le + lem4+ K·12 Arithmetic modulo m m is a positive int. For any int a a mod m 0, 1, 2, ..., m-1 Zm = {0, 1, 2, ..., m-1} +m: atmb = atb (modm) ·m: a·mb= a-b (modm) M = 11 $7 + 9 = 16 \pmod{1} = 5$ 7 · m9 = 63 (mod 11) = 8 **Properties of +**_m and *_m operations if a and $b \in \mathbb{Z}_m$, $C \neq_m b \in \mathbb{Z}_m$ **Closure:** and ambezm Associativity: (a + mb) + mc = a + m(b + mc) = a + mb + mc $(a \cdot mb) \cdot m(= a \cdot m(b \cdot mc) = a \cdot mb \cdot mc$

Commutativity: atmb = btma amb = b ma **Distributivity:** $(a + b) \cdot c = (a \cdot c) + (b \cdot c)$ A + 0 = aIdentity elements 0 and 1 such that a ... 1 = a **Additive inverse:** if a E Zm, then there exists a be Zm such that a + b = 0, a = 7, m=11, then 11-7=4 is the additive inverse of a multiplicative inverse if m is prime

Applications of congruences

1. Hash functions Store several records with keys, in m locations. # Keys < m key vange >> m Solution: Let m=100 banner id = 00975681 ____ loc, 00975681 mod 100 = loc, 81 banner id 00687581 -> loc. 81 2. pseudorandom numbers (RN) - a seguence of deterministically generated numbers Juch that (1) consecutive numbers appear to be unrelated (ii) repetition cycle is large RN Sequence Xo, X1, ---, XO, X1, --- $X_{n+1} = (a \cdot x_n + c) \mod a$, c constants, (n+1)th nth (Linear congruential method)

Problem 7, Section 4.5 ×n+1 = 3 ×n (mod 11). Here a=3, (=0, m=1) X2 = 3. X1 (mod 11) X3 = 3x7 mod 11 = 10 $X_0 = 2$ = 3.6 (mod 11) ×4=3×10mod 11 = 8 X, = 3. Xo (mod ") -18 (mad 11) ×5= 3×8 mod 11= 2=×0 = 3.2 mod (1 =7 =6 The pseudorandom sequence is 2, 6, 7, 10, 8, 2, ... 3. Encryption/decryption Alphabet A, ..., Z. Given a memage, replace each letter by the third letter to the right. for x, y and 2, use A, B and C, respectively. B C D E F S H (A JK 2 12 N O P Q R S T U V W X Y Z 13 19 20 21 22 23 24 25 clear Text UTSA cipher text $f(p) = p + 3 \pmod{26}$ Encrypt: replace each letter by the third letter to position of the letter position of a the left to get clear text. that replaces the letter in the For A, B, C use X, Y, Z, letter in the clear text alphabet respecitively. 5(9)= 8-3 (mod 26) Decrypt: position of the letter position of that replaces the a letter in letter in the cyphier text the alphabet