

Moving Average Frequency Reduction for Low Power in Hard Real Time Systems

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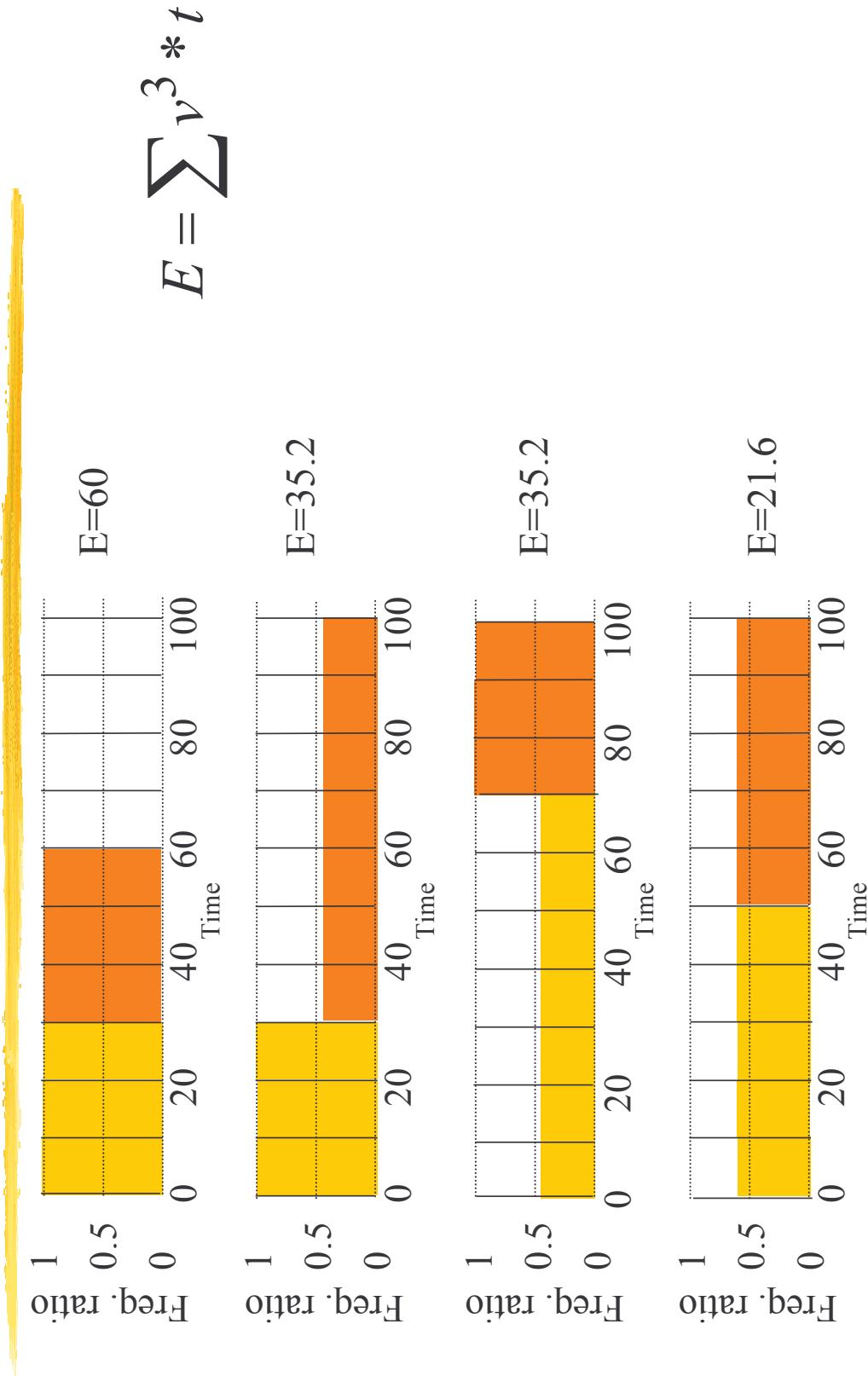
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Low Power in Real Time OS

- Mandatory:
 - Guaranteeing all deadlines $\sum \frac{C_i}{T_i} \leq 1$
 - Hardware allows DVS
- **How to do it?**



Different Policies



When the processor can reduce the speed?

- Periodic tasks in Dual Priority scheduling (DP):
 - 3 queues:
 - URG with a RMA priority policy assignment
 - ARQ with a FIFO policy for aperiodic tasks
 - LRQ with any priority policy assignment (Least Promotion First)
 - Tasks promotes whenever they need to hurry
 - Low Power algorithm:
 - No tasks in the URQ \Rightarrow **minimum frequency**
 - 1 task in the URQ \Rightarrow reduce the frequency meeting the deadlines
 - > 1 task in the URQ \Rightarrow **maximum frequency**



Average frequency reduction policy

First approach:

$$BU = \sum_{\forall \tau_i} \frac{\alpha C_i}{T_i}$$

Where α the minimum increment that will cause at least one task to miss its deadline

We propose:

Lower Bound for frequency reduction = $1/\alpha$



Enhanced Power Low Dual Priority Scheduling

- How it works:
 - Execute the highest priority task in the highest priority queue with a fixed priority assignment (original DP)
 - Execute periodic tasks at **average frequency reduction** provided the time constraints impose by the HRT systems
 - Adjust the clock speed along with voltage supply to reduce power consumption (DVS)



Clock frequency in EPLDP (1)

- Run queues empties

Time power down = $ta_k - \text{wakeup delay}$

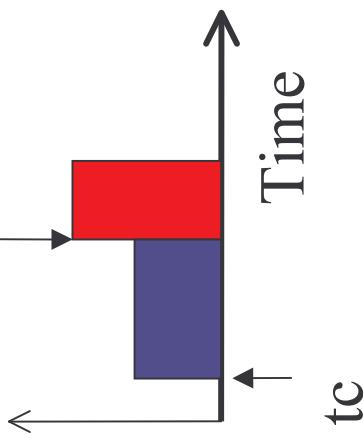
- There is only one task in the URQ

$$Freq\ ratio = \max\left(\frac{\min(tp_k - tc, remaining(C_i))}{\min(tp_k, td_i) - tc}, \frac{U_{rem}}{BU}\right)$$

- More than one task in the URQ

Freq ratio = 1

// Maximum clock speed



Clock frequency in EPLDP (2)

- there are an active LRQ task and URQ empty

if $tp_k < tp_i$ then

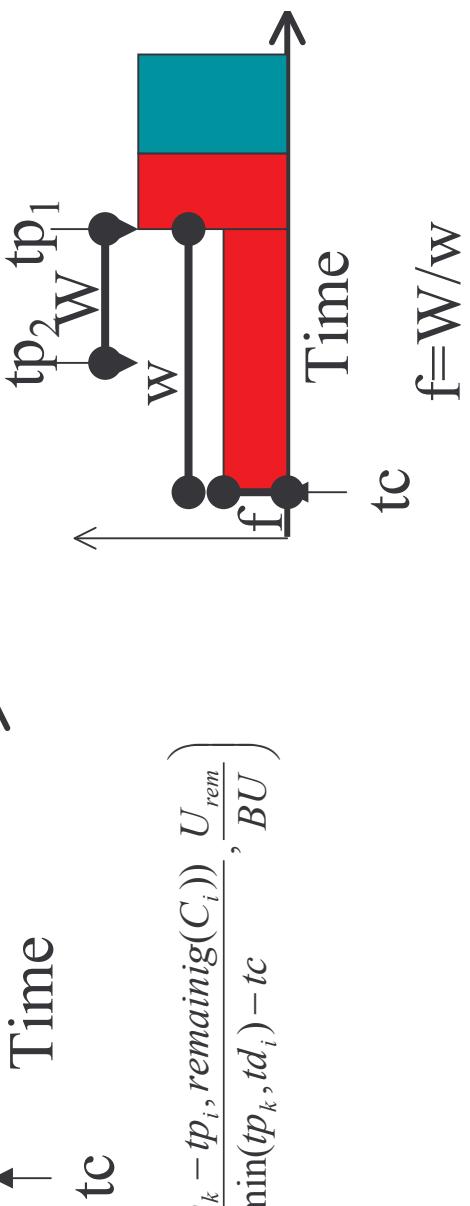
$$tp_k = tp_1$$

$$tp_i = tp_2 / \text{minimum efficient speed}$$

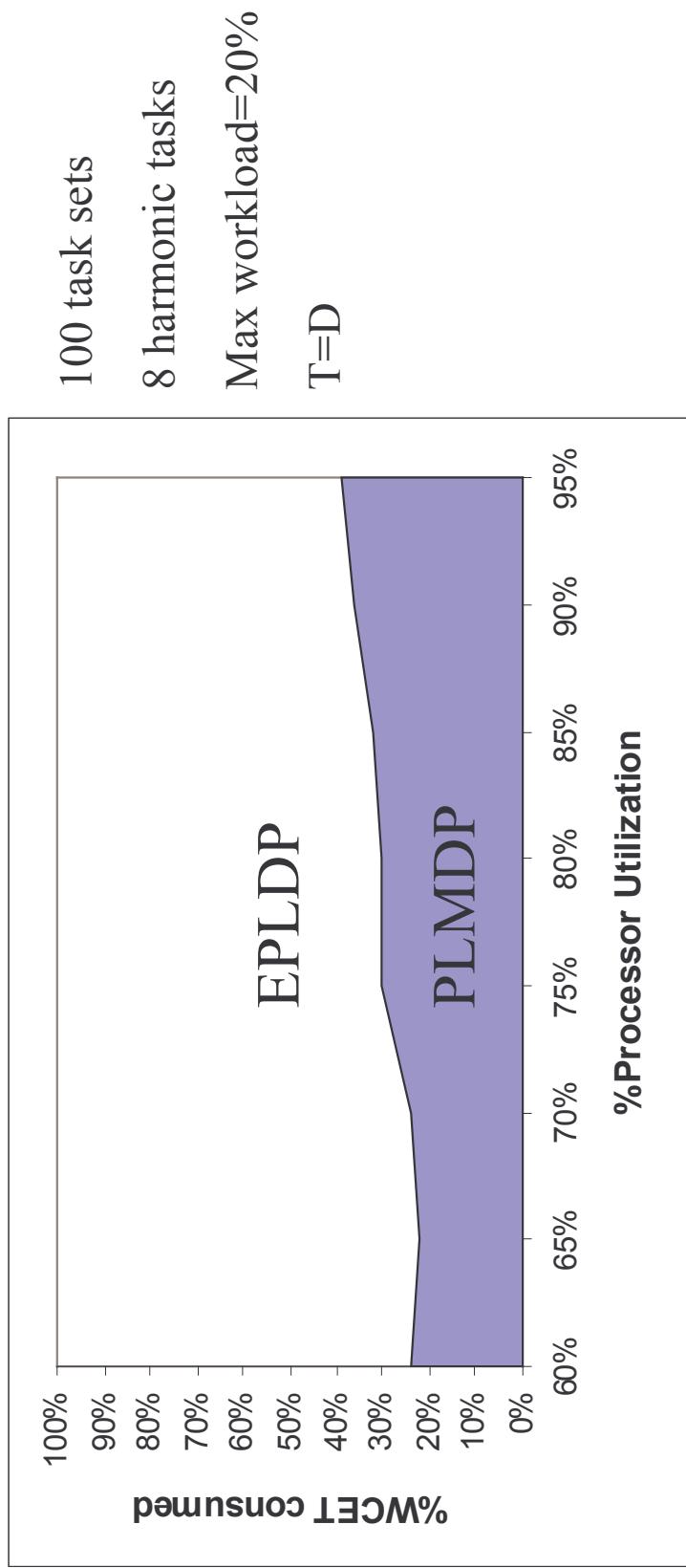
$$Freq\ ratio = \frac{U_{rem}}{BU}$$

else

$$Freq\ ratio = \max \left(\frac{\min(tp_k - tp_i, remaining(C_i))}{\min(tp_k, td_i) - tc}, \frac{U_{rem}}{BU} \right)$$



Performance Transition EPLDP-PLMDP



Moving Average

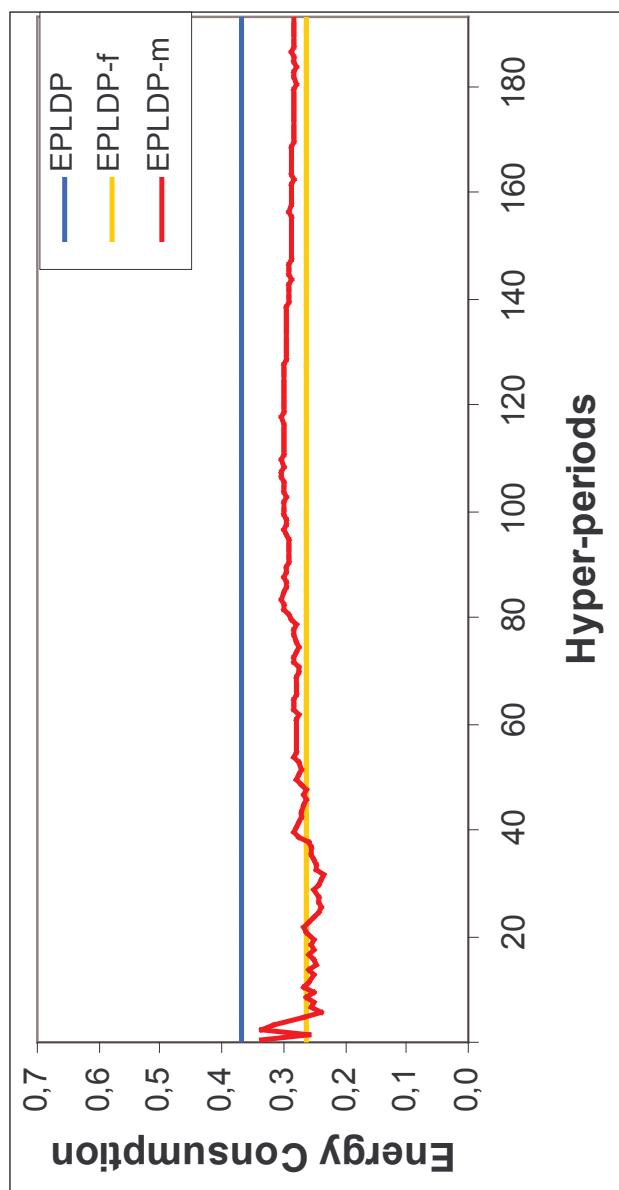
WCET estimation: Moving Average

```
L1       $U_0 = \sum_{\forall \tau} C_i$      $EU = U_0$ 
L2       $i=1;$ 
L3      while real_time_application_not_finished do
L3          execute hyper-period  $i$  with  $EU$  and update  $U_i$ ;
L4       $\overline{U} = \frac{(\overline{U} * i) + U_i}{i+1}$          $EU = U_i - \overline{U}$ 
L5      enddo
```



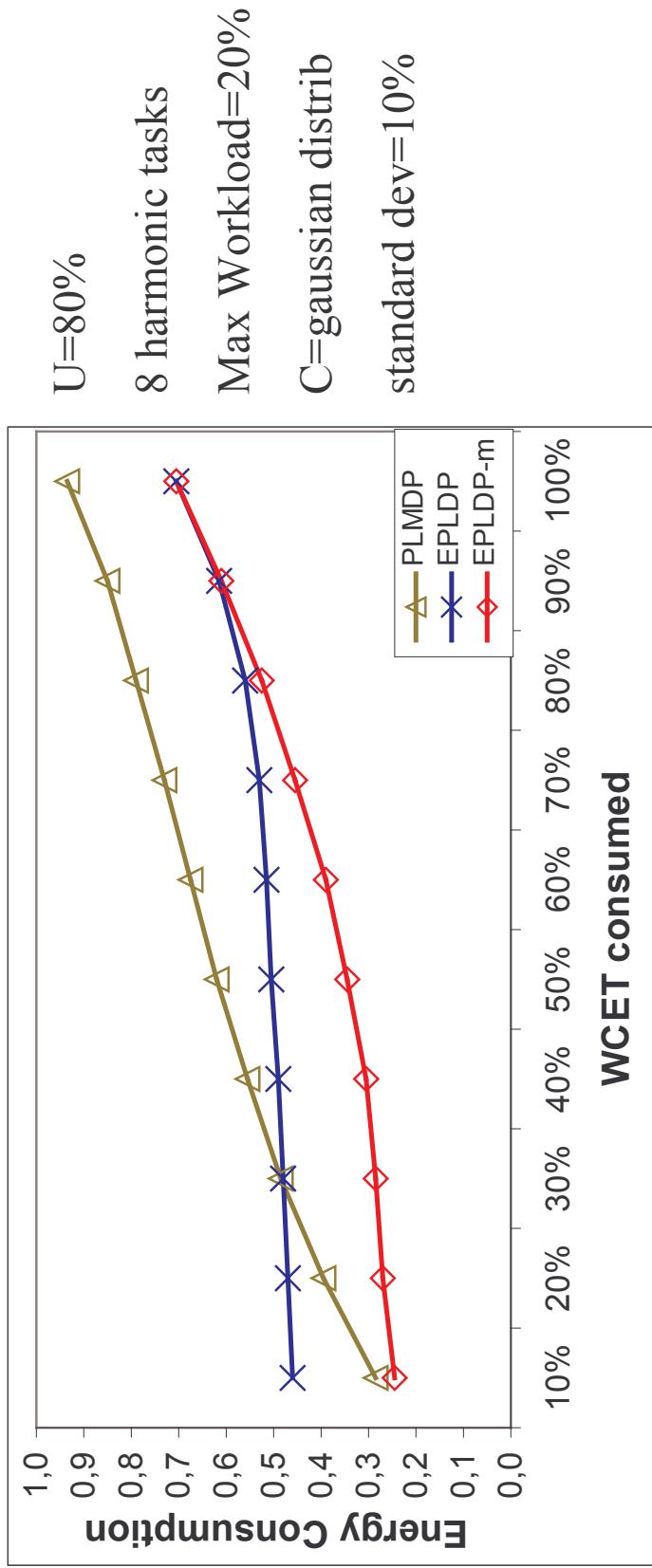
Moving Average

Moving average results

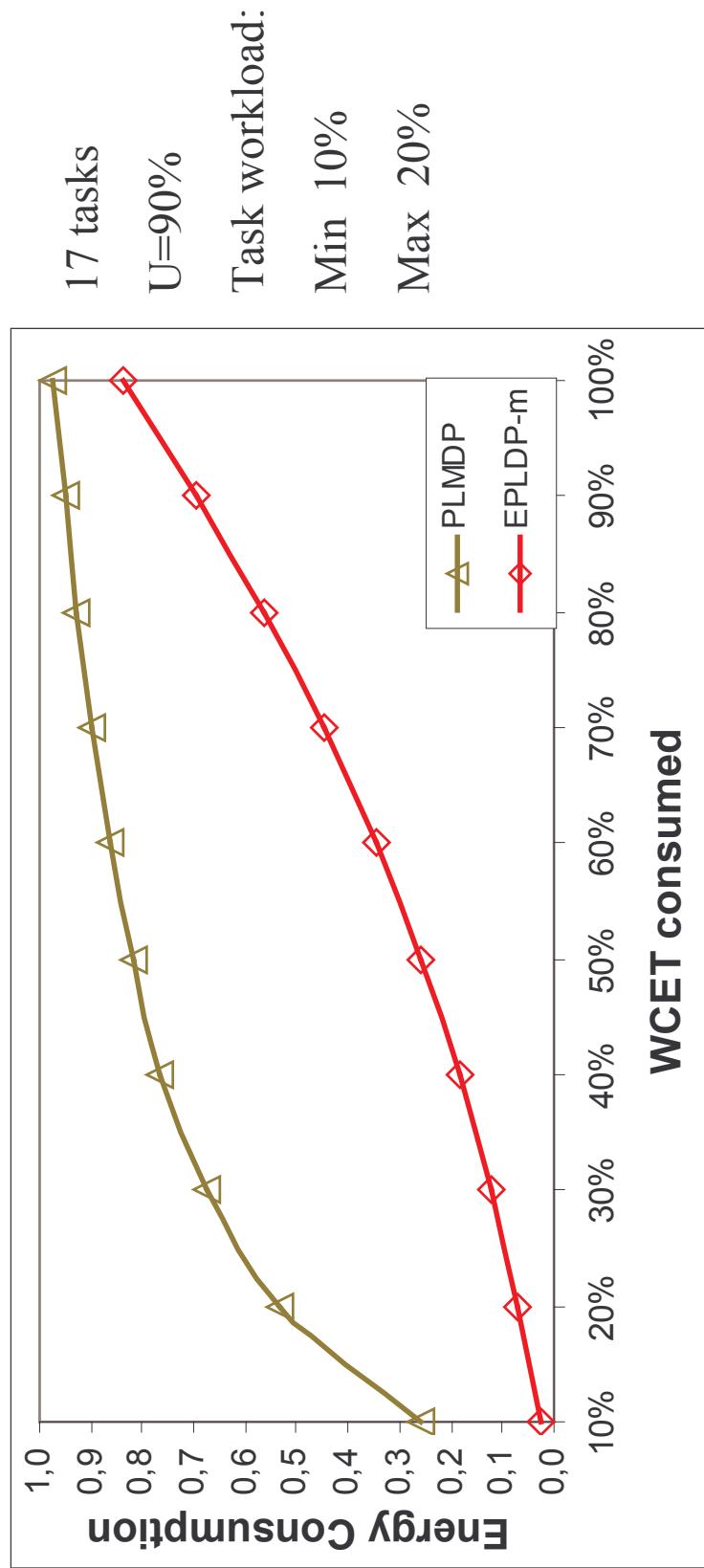


Results

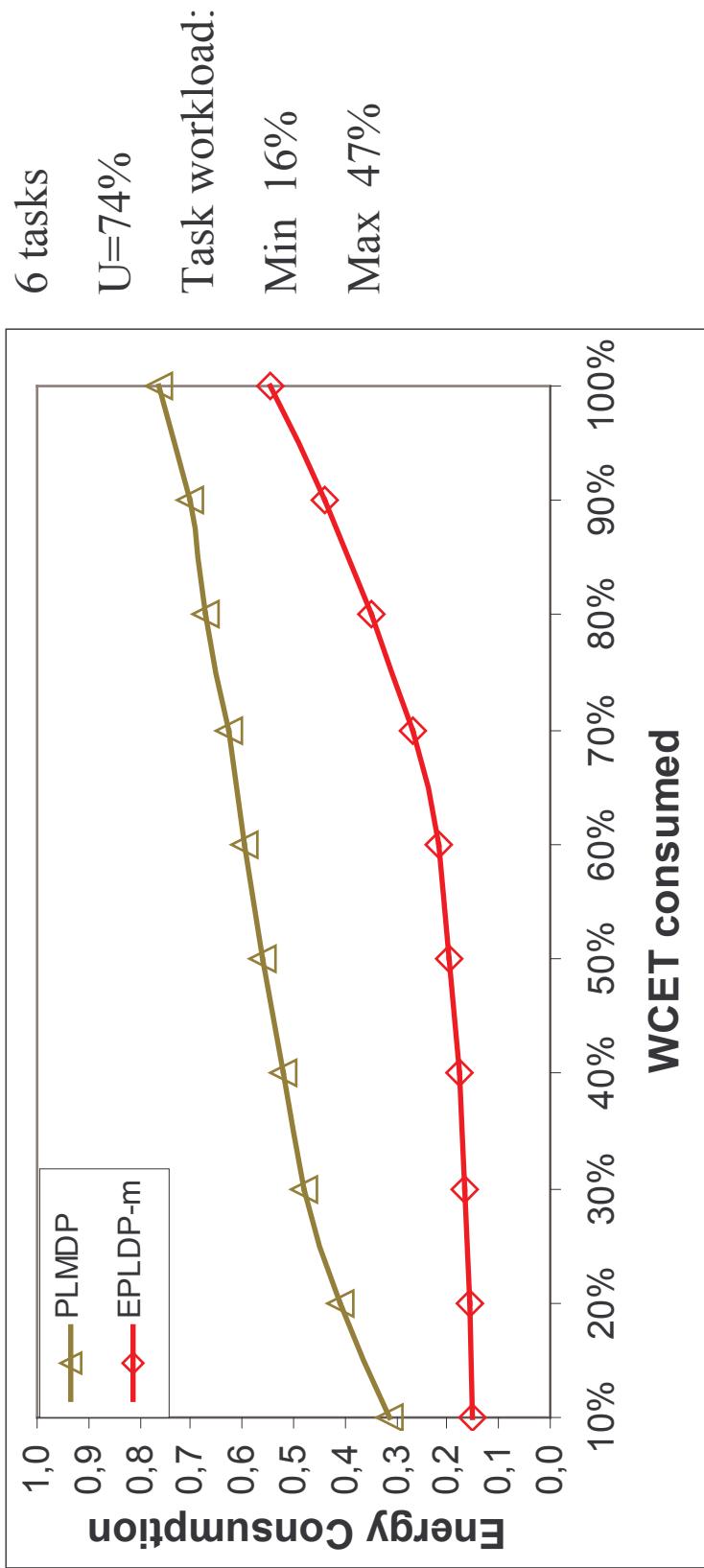
Comparative synthetic benchmark



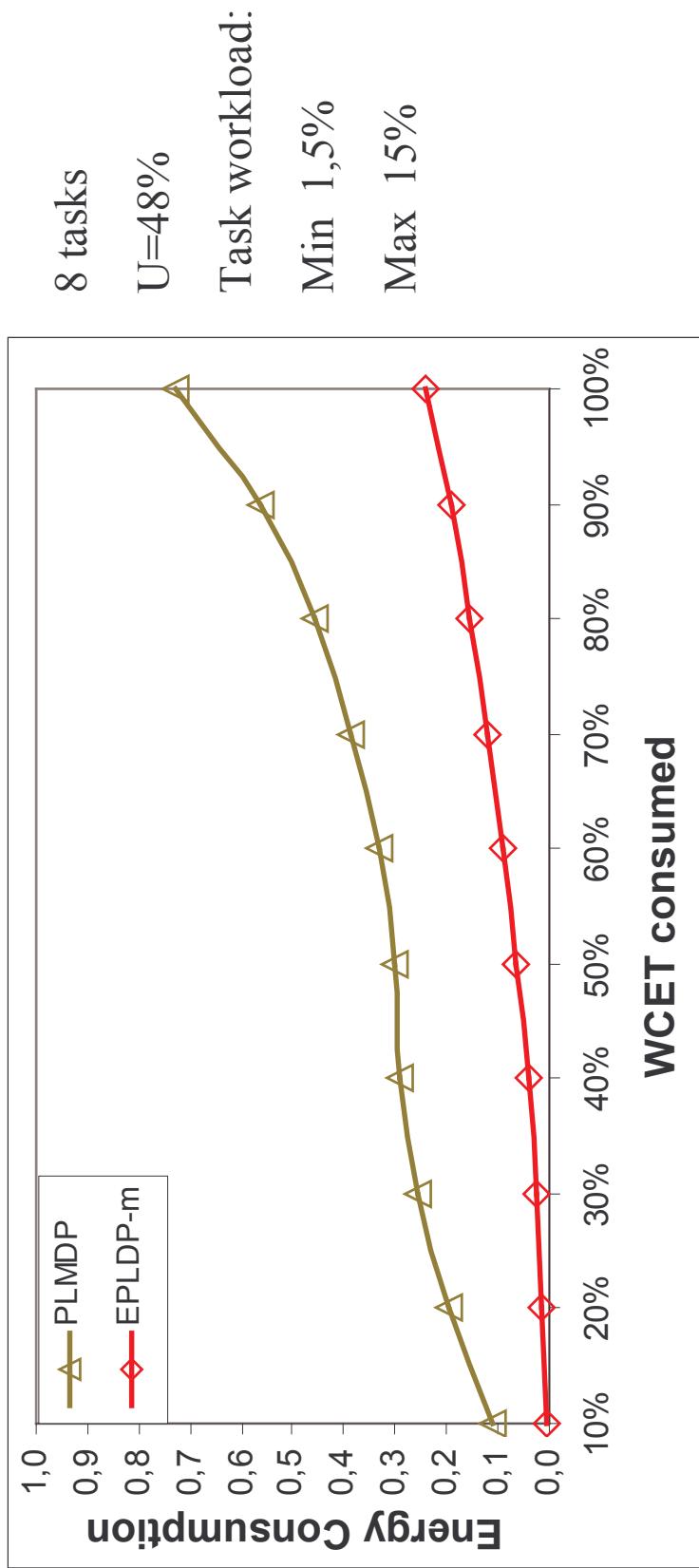
Avionics task set benchmark



INS task set benchmark



CNC task set benchmark



Mean Energy Consumption Comparison

	Avionics	INS	CNC
• LPFPS*	97%	63%	91%
• PLMDP	77%	56%	30%
• EPLDP	79%	41%	22%
• EPLDP-m	35%	27%	9%

* Y. Shin and K. Choi, "Power conscious Fixed Priority scheduling in hard real-time systems"
DAC99, New Orleans, Louisiana, ACM 1-58113-7/99/06, 1999.



Conclusions

- The Enhanced Power Low Dual Priority scheduling obtain significant energy savings without missing the tasks deadlines
- The EPMDP improves PLMDP whenever a moving average is used to determine the real % of WCET used