CS 2213 Advanced Programming Ch 16 – Graphs (Networks) Definitions and Graph Representations

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Thanks to Eric S. Roberts, the author of our textbook, for providing some slides/figures/programs. I also used some materials from other textbooks (specifically, The Algorithm Design Manual by **Steven Skiena** and his slides at <u>http://www.cs.sunysb.edu/skiena</u>. So I also thank to **Steven Skiena** for presentations on the Internet).

Disclaimer

- You are highly recommended to read the chapter 16 from our textbook, which provides a general interface-based design for graph abstraction using several other data structures (e.g., list, set) covered in previous chapters.
- Instead of this general abstract design, we will directly implement concrete graph structures and operations. So even though the key concepts are the same, the implementation details will be different than the textbook.

Objectives

- To appreciate the **conceptual** structure of a graph and its **applications**
- To learn basic graph theory terminology/notation
- To learn underlying representations for graphs (adjacency matrix, adjacency list)
- To understand and be able to apply basic graph algorithms (topological sort, depth-first search, breath-first search, Dijkstra's shortest path etc.)
- To provide knowledge and skills to students so that they can comfortably use graph structures and algorithms (or develop new ones) in their research or workplace.

Structure of a graph

A **graph** (or network) is a mathematical object consisting of **nodes** (also called vertices) and **edges** (also called arcs, links).



The airline graph represented by G(V,E) consists of the following sets.

- V = { Atlanta, Boston, Chicago, Dallas, Denver, Los Angeles, New York, Portland, San Francisco, Seattle }
- $\begin{array}{l} {\it E} = \{ \mbox{ Atlanta} \leftrightarrow \mbox{ Dallas, Atlanta} \leftrightarrow \mbox{ New York,} \\ \mbox{ Boston} \leftrightarrow \mbox{ New York, Boston} \leftrightarrow \mbox{ Seattle, Chicago} \leftrightarrow \mbox{ Denver,} \\ \mbox{ Dallas} \leftrightarrow \mbox{ Denver, Dallas} \leftrightarrow \mbox{ Los Angeles, Dallas} \leftrightarrow \mbox{ San Francisco,} \\ \mbox{ Denver} \leftrightarrow \mbox{ San Francisco, Portland} \leftrightarrow \mbox{ San Francisco, Portland} \leftrightarrow \mbox{ Seattle} \} \end{array}$

Applications of graphs

Graphs (networks) are everywhere!

Physical Networks

- Transportation networks
- Power transmission networks
- Telephone/computer networks

What are the common things in all these problems?

 We move some entity {car, people, electricity, voice, data} from one node {intersection, city, generator, telephone, computer} to another through underlying links {roads, power lines, communication links}.

But, we want to move entities from one point to another in an efficient way (what do we mean?)

Application: Shortest path

Assume each link has cost (e.g., distance).



What is the best way to get from A to B as <u>cheaply</u> as possible? How can we find it?

Finding Shortest path

- Enumerate all possible solutions
 - Measure the "cost" of each alternative solution
 - Select the one with min cost
- From the perspective of pure mathematics, the above problem is trivial to solve! BUT
 - The number of paths could be extremely large

2ⁿ possible solutions, what happens when n=100
So, we need efficient graph algorithms.

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Basic Graph Notation

Basic Graph Notation

Network or Graph

- G = (N, A) or G(V, E)
 - $V = \{ v_1, v_2, ..., v_n \}$





node/vertex set

- arc/link/edge set
- Link/edge connecting v_p to v_q is also denoted by (p,q)

Incidence and adjacency:

- e_i and e_i are adjacent links
- v_p and v_q are adjacent nodes
- e_i and e_i are incident to node v_a

Labeled vs. Unlabeled Graphs

 Each vertex is assigned a unique name or identifier to distinguish it from other vertices.



 An important graph problem is *isomorphism testing*, determining whether the topological structure of two graphs are in fact identical if we ignore any labels.

Weighted vs. Unweighted Graphs

- In unweighted graphs, there is no cost distinction between various edges and vertices.
- In weighted graphs, we can associate weights/labels/attributes with nodes or links,
 - for node p: **w(p)**, or w_p
 - for link (p,q): c(p,q) or c_{pq}





Simple vs. Non-simple Graphs

- A *self-loop* is an edge (x,x) involving only one node
- An edge (x,y) is a *multi-edge* if it occurs more than once in the graph.



 Any graph which avoids these structures (loops, multi edges) is called *simple*.

Directed and Undirected graphs:

- Directed network:
 - Flow is in one direction, e.g., from v_p to v_q
- Undirected network:
 - Flow can be in any direction with the same attributes.
- If link attributes are different, we can use directed network model



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An *acyclic* graph does not contain any cycles.



Directed Acyclic Graphs are called DAGs.

They arise naturally in scheduling problems, where a directed edge (x, y) indicates that x must occur before y.

- **Degree**: the number of incident to a node
 - deg(v_q) is 3
 - If Directed, in-degree, out-degree



- Path: a distinct sequence of nodes
 - Undirected vs. directed
- Walk: a sequence of nodes
- Cycle: a distinct sequence of nodes except the first and last node

Connected graph

- There is at least one path between every pair of nodes.
- Connected components
- Incase of directed graphs
 - Strongly connected
 - Weakly connected

Tree

A connected graph with no cycles

Sparse vs. Dense Graphs

- Graphs are *sparse* when only a small fraction of the possible number of vertex pairs actually have edges defined between them.
- Graphs are *dense* when the number of edges increases
- Complete graph: there is an edge between every pair of nodes
- Dense graphs have a quadratic number of edges while sparse graphs are linear in size.







Some important nodes

- **Source node**: (or starting point) a node where flow originates
- Destination/Sink: (or termination point) a node where flow ends
- Supply node: a node where supply for flow is available
- Demand node: a node where non-zero demand for flow exists



Example: The Friendship Graph



- Consider a graph where the vertices are people, and there is an edge between two people if and only if they are friends.
- What questions might we ask about the friendship graph?
 - If I am your friend, does that mean you are my friend? (Undirected)
 - Am I my own friend? (self-loop, simple)
 - Am I linked by some chain of friends to the President? (path)
 - How close is my link to the President? (shortest path)
 - Is there a path of friends between any two people? (strongly connected)
 - Who has the most friends? (degree)
 - Is there a clique? a clique in an <u>undirected graph</u> is a subset of its <u>vertices</u> such that every two vertices in the subset are connected by an edge.

Data Structures for Graphs

Suppose the graph G = (V, E) contains |V|=n vertices and |E|=m edges.



How can we represent this graph in C?



- We can represent G(V, E) using an nxn matrix M, where element M[i][j] will be
 - I if (i,j) is an edge of G, and
 - 0 if it isn't.
- Is this a good representation
 - for dense graphs
 - for sparse graphs,
- Can we save space,
 - if the graph is undirected?
 - if the graph is sparse?



Each element M[i][j] might be a structure to represent other attributes of link (i,j)

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Exercise: Adjacency Matrix for the complement graph



Modify this matrix



Adjacency Lists

- We can represent G(V, E) using an n×1 array of **pointers**, where the ith element points to a **linked list** of the edges incident on vertex *i*.
- Is this a good representation
 - for dense graphs
 - for sparse graphs



Adjacency Lists vs. Adjacency Matrices

- Which representation would be better for the followings, also give complexity:
 - Test if (x, y) exists?
 - Find vertex degree?
 - Memory usage for sparse graphs?
 - Memory usage for dense graphs?
 - Edge insertion or deletion?
 - Faster to traverse the graph?
- Better for most problems?

Matrices O(1)Lists O(degree) lists (m+n) vs. (n^2) matrices O(n²) matrices O(1)lists (m+n) vs. (n^2) ists 24

Adjacency List Representation in C

Concrete Graph Structure

#define MAXV 100

typedef struct edgenode {
 int y;
 int w; //weight
 struct edgenode *next;
} edgenodeT;

```
typedef struct {
    edgenodeT *edges[MAXV+1];
    int degree[MAXV+1];
    int nvertices;
    int nedges;
    bool directed;
} graphT;
```

We will assume that we will have at most MAXV many nods in our graphs; but , we can dynamically change graph size. How?

adjacency info x - - - - > y link weight/cost if any pointer to next link

The edges field is an array of pointers to list of edges originating from each node x.
The degree field counts the number of meaningful entries for the given vertex.
An undirected edge (x, y) appears twice in any adjacency-based graph structure, once as y in x's list, and once as x in y's list.

Creating a gra a client progra		oh m	ir) g1	typ e ir ir ir b } gr	edef dgen it deg it nve it nec ool d aphT	struct { odeT *ec gree[MA> ertices; dges; irected; ;	lges[(V+1)	мах ^у]; уg2	V +1] ;
	, [de	gree	e	dges]	\sim		•	
main() // Assume that MAXV is 6		0	?	0	?		v∕ de	gree	e	dges
{ graphT myg1, *myg2;		1	?	1	?		0	?	0	?
		2	?	2	?		1	?	1	?
myg2 = (graphT *) malloc(sizeof(graphT));		3	?	3	?		2	?	2	?
if(myg2==NULL) {		4	?	4	?		3	?	3	?
<pre>printf("no memory");</pre>		5	?	5	?		4	?	4	?
exit(-1); } // myg2 = New(graphT *);		6	?	6	?		5	?	5	?
initialize_graph(&myg1, TRUE);		nverti	ces	?]		6	?	6	?
		nedge	es	?			nverti	ces	?	
initialize_graph(myg2, FALSE);		direct	ed	?			nedge	es	?	
}						-	direct	ed	?	



typedef struct {
 edgenodeT *edges[MAXV+1];
 int degree[MAXV+1];
 int nvertices;
 int nedges;
 bool directed;
} graphT;











Print Graph

typedef struct edgenode {	typedef struct {
int y;	edgenodeT
int w;	*edges[MAXV+1];
struct edgenode *next;	int degree[MAXV+1];
} edgenodeT;	int nvertices;
·	int nedges;
	bool directed;
	} graphT;

print_graph(graphT *g) edges degree 0 0 ? edgenodeT *pe; 2 1 int i; for(i=1; i<=g->nvertices; i++) { 2 3 2 printf("Node %d: ", i); 3 3 3 pe = g->edges[i]; 3 4 4 while(pe){ 5 3 5 printf(" %d %d,", pe->y, pe->w); pe = pe->next; 6 2 6 } nvertices 6 printf("\n"); nedges 8 directed 0

Free Graph

typedef struct edgenode {	typedef struct {				
int y;	edgenodeT *edges[MAXV+1]:				
struct edgenode *next; } edgenodeT;	int degree[MAXV+1];				
	int nvertices;				
	int nedges;				
	bool directed;				
	} graphT;				

free_graph(graphT *g) degree edges 6 0 ? 0 edgenodeT *pe, *olde; 2 1 int i; for(i=1; i<=g->nvertices; i++) { 2 2 3 pe = g->edges[i]; 3 3 3 while(pe){ 3 4 4 olde = pe;5 3 5 pe = pe->next; free(olde); 6 2 6 } 6 nvertices 46 8 nedges free(g); directed 0



Exercise: Graph copy

recitation

- Write a function that creates a new copy of a given graph g using adjacency list or adjacency matrix
- void copy_graph_list(graphT *g, graphT **newg)

Why ** ?

```
graph *copy_graph_list(graphT *g)
{ ....
}
```





MORE EXERCISES

Exercise: Convert graph representation

- Suppose you are given a graph g, which is represented using
 - adjacency list or
 - adjacency matrix,
- Write a function to convert it to a new graph newg using
 - adjacency matrix or
 - adjacency list, respectively...
- See next slide for an example...



convert_list_to_matrix Modify print_graph



Find in-degree and out-degree for each node



